

Introductory Graph Theory

1. Definition of a Graph

Intuitive Definition: A *simple graph* is a collection of vertices (visualized as dots) and edges (visualized as arcs between dots).

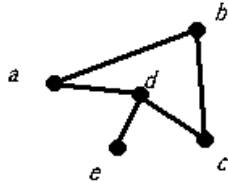


Figure 1.1

Formal Definition: A *simple graph* G with n vertices and m edges consists of a vertex set $V(G) = \{v_1, v_2, \dots, v_n\}$ and edge set $E(G) = \{e_1, e_2, \dots, e_m\}$, where $E(G) \subseteq V(G) \times V(G)$, that is each edge is an unordered pair of vertices.

Definitions: A *vertex* v is *incident to an edge* if v is one of the pair of vertices which determines the edge. The *degree of a vertex* is the number of edges to which it is incident. We denote the degree of a vertex v as $\deg(v)$. Given a graph G on vertex set $V(G) = \{v_1, v_2, \dots, v_n\}$, the *degree sequence of G* is $\deg(v_1), \deg(v_2), \dots, \deg(v_n)$. The maximum value of the degree sequence is the *maximum degree of the graph* and we denote it by $\Delta(G)$. The minimum value of the degree sequence is the *minimum degree of the graph* and we denote it by $\delta(G)$.

2. Some Simple Applications of Graphs

Acquaintance Graph. Suppose that the the people are vertices and that there is an edge between two people if they are acquaintances.

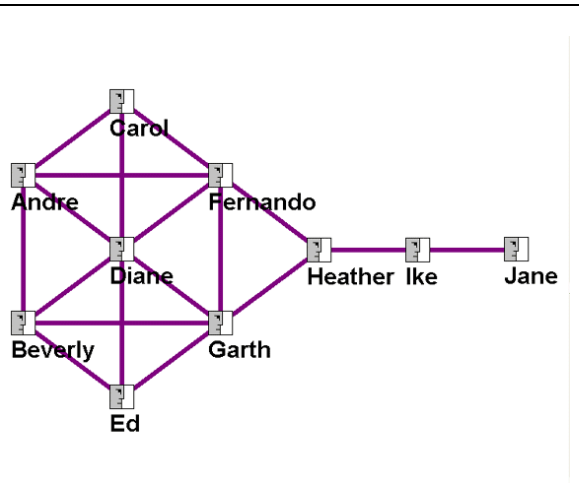
Exercise: Determine the following.

The number of vertices is _____

The maximum degree is _____

The minimum degree is _____

Interpret the maximum and minimum degree.



Network Graph. Suppose that the computers are vertices and that there is an edge between two vertices if they have direct communication.

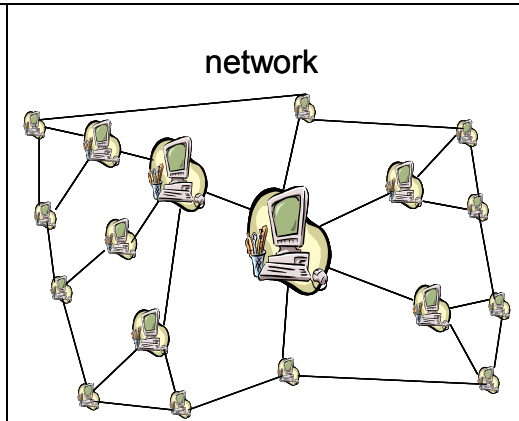
Exercise: Determine the following.

The number of vertices is _____

The maximum degree is _____

The minimum degree is _____

Interpret the maximum and minimum degree.

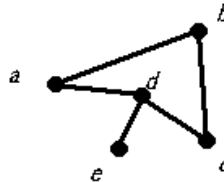


3. Some Special Graphs and More Graph Concepts and Invariants

A *graph concept* is a property of the graph. A *graph invariant* is a numeric value associated with the graph, usually independent of the way the graph is drawn.

Definition: Let G be a graph on vertex set V and edge set E . For any two vertices, say u and v , if (u, v) is an edge we say that u is *adjacent to* v . To indicate that two vertices u and v are adjacent we use the notation $u \sim v$.

Example: Since (a, b) is an edge
 $a \sim b$.



Definition: The *empty graph on n vertices* (also called the null graph on n vertices) is the graph on n vertices with no edges.

Definition: The *complete graph on n vertices* is the graph on n vertices in which every two vertices are adjacent. We use the notation K_n to denote the complete graph on n vertices.

Definition: Let G be a graph on vertex set V and edge set E . We define the **complement graph of G** , denoted \overline{G} , as a graph on the the same vertex set V in which two vertices adjacent in \overline{G} if and only if they are not adjacent in G .

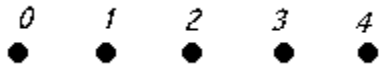
It is easily seen that the empty graph on n vertices is denoted by $\overline{K_n}$.

Exercise : Draw \overline{G} if G is the graph drawn below.

Exercise : Find a graph G (on four vertices) for which G and \overline{G} can be drawn so they *appear* the same; formally this should be so that $G \simeq \overline{G}$. The symbol \simeq is used to indicate that the two graphs are essentially the same except for the way the vertices are labeled (formally the symbol \simeq is read as *isomorphic to*, which we will investigate soon.)

Exercise : Suppose we are given a graph on vertex set $\{0, 1, 2, 3, 4\}$ and that edges of this graph are determined by the following rule:

$\forall x, y \in \{0, 1, 2, 3, 4\} \quad x \sim y \leftrightarrow |x - y| = 1$. Draw the edges subject to the given rule in the following diagram.



Definition: A graph is called a **Path on n vertices** if the vertices can be labeled with elements of $\{0, 1, 2, \dots, n-1\}$ so that the edge set is $\{(i, i+1) \mid i \in \{0, 1, 2, \dots, n-1\}\}$. Such a graph is denoted by P_n .

Exercise: Verify that this graph in the figure is P_6 .



Definition: A graph is called a *cycle on n vertices* ($n \geq 3$) if the vertices can be labeled with elements of $\{0, 1, 2, \dots, n-1\}$ so that the edge set is $\{(i, i+1) \mid i \in \{0, 1, 2, \dots, n-1\}\} \cup \{(0, n-1)\}$. Such a graph is denoted by C_n .

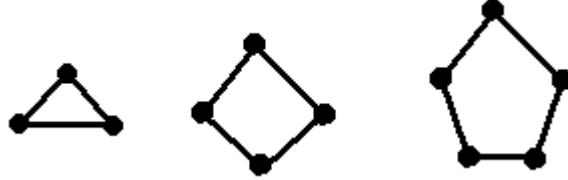


Figure 1.3. C_3 , C_4 , and C_5

Definition: A graph is called a *Wheel on n vertices* if the vertices can be labeled with elements of $\{0, 1, 2, \dots, n-1\}$ so that the vertices $\{1, 2, \dots, n-1\}$ determine a cycle on $n-1$ vertices and vertex 0 is adjacent to each of the vertices in $\{1, 2, \dots, n-1\}$. Such a graph is denoted by W_n .

Exercise : Verify that the following graph is W_6 .



Definition. The **n-Cube graph**, denoted Q_n , is the graph that has vertices representing the 2^n bit strings of length n . Two vertices are adjacent if and only if the bit strings that they represent differ in exactly one position.

Exercise : Draw the edges of the 3-cube in the next figure.

